

Markov Chain Quantal-Response

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Abstract

Real-world players in games are susceptible to making errors. In the quantal-response model, the players perceive utilities that include random errors. Here, the players' responses are probabilistic, with players having a higher likelihood of choosing better responses than worse responses. We study the dynamics of the quantal-response model by analyzing the stationary distribution of a Markov chain. The probability of transitioning from one strategy profile to another is determined by quantal-responses. The stationary distribution is a joint probability distribution on the players' actions. We consider a game in which the players are divided into two groups, and the utility of a player in one group depends only on the actions of the players in the other group. We find that if we allow all the players in such a game to simultaneously revise their responses,

* listed in alphabetical order

the stationary distribution is the product of the marginal distributions of each group.

1 Definitions and Notation

1.1 Quantal-Response Normal-Form Game

Let $G = (N, S, u)$ be a normal-form game. $N = \{1, 2, \dots, n\}$ is the set of players. For each player $i \in N$, there is a strategy set $S_i = \{s_{i1}, s_{i2}, \dots, s_{iJ_i}\}$ of J_i pure strategies. Each player $i \in N$ also has a payoff function $u_i : S \rightarrow \mathbb{R}$, where $S = \prod_{i \in N} S_i$.

We denote $\Delta = \prod_{i \in N} \Delta_i$ as the set of mixed strategy profiles, where Δ_i is the set of probability measures on S_i . For all $p_i \in \Delta_i$, $p_i : S_i \rightarrow \mathbb{R}$, where $\sum_{s_{ij} \in S_i} p_i(s_{ij}) = 1$, and $p_i(s_{ij}) \geq 0, \forall s_{ij} \in S_i$. To simplify notation, we denote $p_{ij} = p_i(s_{ij})$.

The payoff function of each player $i \in N$ is $u_i(p) = \sum_{s \in S} p(s)u_i(s)$, where $p(s) = \prod_{i \in N} p_i(s_i)$. We denote $\zeta_i = \mathbb{R}^{J_i}$ as the space of possible payoffs for strategies that player i could adopt. Let $\zeta = \prod_{i=1}^n \zeta_i$. We define $\bar{u} : \Delta \rightarrow \zeta$, by

$$\bar{u}(p) = (\bar{u}_1(p), \dots, \bar{u}_n(p)), \text{ where } \bar{u}_i(p) \in \mathbb{R}^{J_i}$$

and where

$$[\bar{u}_i(p)]_j = \bar{u}_{ij}(p) = u_i(s_{ij}, p_{-i}).$$

We consider the scenario in which each player's utility for each action is subject to random error. For each player $i \in N$, $\forall j \in \{1, \dots, J_i\}, \forall p \in \Delta$,

$$\hat{u}_{ij}(p) = \bar{u}_{ij}(p) + \varepsilon_{ij}.$$

Player i 's error vector is defined as $\varepsilon_i = (\varepsilon_{i1}, \varepsilon_{i2}, \dots, \varepsilon_{iJ_i})$. The errors ε_i are distributed according to joint density $f_i(\varepsilon_i)$. We say that $f = (f_1, \dots, f_n)$ is admissible if the marginal distribution of each f_i exists for each ε_{ij} and $\mathbb{E}[\varepsilon_i] = \mathbf{0}$.

Following the notation in [2], given $p \in \Delta$, for all $\bar{u}(p) = (\bar{u}_1(p), \dots, \bar{u}_n(p))$, where $\bar{u}_i(p) \in \mathbb{R}^{J_i}, \forall i$, the ij-response set $R_{ij}(\bar{u}_i(p)) \subseteq \mathbb{R}^{J_i}$ is defined as

$$R_{ij}(\bar{u}_i(p)) = \{\varepsilon_i \in \mathbb{R}^{J_i} \mid \bar{u}_{ij}(p) + \varepsilon_{ij} \geq \bar{u}_{ik}(p) + \varepsilon_{ik}, \forall k = 1, \dots, J_i\}.$$

In other words, the ij-response set $R_{ij}(\bar{u}_i(p))$ is the set of all error vectors ε_i that make action j the dominant strategy for player i when all other players perform strategy profile p . The probability that player $i \in N$ will select strategy $j \in S_i$ when given $\bar{u}_i(p) \in \mathbb{R}^{J_i}$ is the quantal-response function $\sigma_{ij}(\bar{u}_i(p))$, which is defined as

$$\sigma_{ij}(\bar{u}_i(p)) = \int_{R_{ij}(\bar{u}_i(p))} f(\varepsilon) d\varepsilon.$$

1.2 Logit-Response (note: not used in Theorem)

A particularly interesting case of the quantal-response function is the logistic quantal-response function. Recall $f = (f_1, \dots, f_n)$, where the errors $\varepsilon_i \in \mathbb{R}^{J_i}$ are distributed according to joint density $f_i(\varepsilon_i)$. For some $\beta \geq 0$, the logistic quantal-response function gives the optimal choice behavior when f_i , $\forall i \in N$, is an extreme value distribution with the cumulative distribution function $F_i(\varepsilon_{ij}) = e^{-e^{-\beta\varepsilon_{ij}-\gamma}}$ and the errors ε_{ij} are independent. Specifically, for some $\bar{u}_i(x) \in \mathbb{R}^{J_i}$, the logistic quantal-response function is

$$\sigma_{ij}(\bar{u}_i(x)) = \frac{e^{\beta u_i(s_{ij}, x-i)}}{\sum_{k=1}^{J_i} e^{\beta u_i(s_{ik}, x-i)}}.$$

1.3 Markov Chain

A Markov chain is a discrete-time stochastic process defined by a sequence of random variables X_1, X_2, \dots, X_n that satisfies the Markov property

$$\mathbb{P}(X_{i+1} = x | X_1 = x_1, X_2 = x_2, \dots, X_n = x_n) = \mathbb{P}(X_{i+1} = x | X_i = x_i),$$

where $\mathbb{P}(X_1 = x_1, X_2 = x_2, \dots, X_n = x_n) > 0$. The random variables X_1, X_2, \dots, X_n take values on the countable state space, which we define as the set of action profiles S in Section 1.1.

The initial state X_1 is chosen at random according to a distribution λ on S , which assigns probability λ_i to state i via

$$\mathbb{P}(X_1 = i) = \lambda_i.$$

The Markov chain has an associated stochastic transition matrix $\mathcal{P} = [p_{ij}]$, where $i, j \in S$. Here, p_{ij} is the probability of state X_k being action profile i and transitioning to action profile j . That is to say $p_{ij} = \mathbb{P}(X_{k+1} = j | X_k = i)$. For each $i \in S$, $\sum_{j \in S} p_{ij} = 1$. Also, $p_{ij} \geq 0$, $\forall p_{ij}$.

1.4 Transition Matrix of the Markov Chain

We consider a few cases of the transition matrix \mathcal{P} that are differentiated by which players receive revision opportunities in a period. Just as in [1], we will define learning dynamics as a behavioral rule (such as logit responding) and a revision process.

1.4.1 Asynchronous Learning (note: not used in Theorem)

We first consider asynchronous learning, in which exactly one player receives a revision opportunity in a period. For any player $i \in N$, we define $m_i(x, y)$ as the likelihood of player i receiving a revision opportunity and changing the current action profile x into the action profile y . In other words, we have

$$m_i(x, y) = \begin{cases} \sigma_{iy_i}(\bar{u}_i(x)) & \text{if } x_{-i} = y_{-i} \\ 0 & \text{if } x_{-i} \neq y_{-i} \end{cases}, \quad (1)$$

where x_{-i} and y_{-i} are the action profiles without the action of player i and y_i is the action of player i in strategy profile y . Then, transition matrix \mathcal{P} is defined as $\mathcal{P} = [p_{xy}]$, where $p_{xy} = \frac{1}{n} \sum_{i=1}^n m_i(x, y)$.

1.4.2 Simultaneous Learning

We also consider the case in which all players receive revision opportunities in every period. Here, the transition matrix \mathcal{P} is defined as $\mathcal{P} = [p_{xy}]$, where $p_{xy} = \prod_{i \in N} \sigma_{iy_i}(\bar{u}_i(x))$.

1.5 Update Function Representation of Markov Chain (note: not used in Theorem)

An equivalent representation of the Markov chain replaces the probabilistic transitions of a Markov chain with a deterministic update function that takes as input realizations of a uniform random variable. Let $U_n, n \geq 1$ be a sequence of independent and identically distributed random variables, in which $U_n \sim \text{Uniform}[0, 1], \forall n \geq 1$. For clarity, let us enumerate the elements in the state space S from 1 to $|S|$. Consider the piecewise function $\phi : S \times [0, 1] \rightarrow S$, such that

$$\phi(i, U_n) = \begin{cases} 1 & \text{if } U_n \in [0, p_{i1}] \\ 2 & \text{if } U_n \in (p_{i1}, p_{i1} + p_{i2}] \\ \dots & \\ j & \text{if } U_n \in \left(\sum_{t=1}^{j-1} p_{it}, \sum_{t=1}^j p_{it} \right] \\ \dots & \\ |S| & \text{if } U_n \in \left(\sum_{t=1}^{|S|-1} p_{it}, 1 \right]. \end{cases}$$

Starting in state $X_n, \forall n \geq 1$, the next state X_{n+1} is determined by the following

$$X_{n+1} = \phi(X_n, U_n).$$

1.6 Graphical Games

Consider the game $G = (V, S, u)$ with $V = \{1, \dots, n\}$, where $n = |V|$ players. Also, consider the simple, undirected graph $\mathcal{G} = (V, E)$ with V as the set of vertices and E as the set of edges. We define the game G as a graphical game on \mathcal{G} if each player's utility is a function of its neighbors on \mathcal{G} . In other words, G is a graphical game on \mathcal{G} if for all $i, j \in V$, if $i \neq j$ and $(i, j) \notin E$, then

$\forall a \in S$, $u_i(a) = u_i(a^j)$ where action profile a^j is the same as a in all but the j^{th} position.

2 Theorem

Let $\mathcal{G} = (V, E)$ be a bipartite graph with partitions \mathcal{G}_0 and \mathcal{G}_1 . Let $G = (V, S, u)$ be a graphical game on \mathcal{G} . If the quantal-response dynamics exhibit simultaneous learning, then the stationary distribution ($\pi = \lim_n \lambda \mathcal{P}^n$, for arbitrary initial distribution λ) of the quantal-response dynamics is a product measure. Specifically,

$$\pi \mathcal{P} = \pi,$$

such that

$$\pi = \pi_0 \times \pi_1,$$

where π_0 is a measure on $\mathcal{S}_0 = \prod_{k \in \mathcal{G}_0} S_k$ and π_1 is a measure on $\mathcal{S}_1 = \prod_{k \in \mathcal{G}_1} S_k$.

2.1 Proof

In the bipartite graph G , we know that the utilities of players in one partition depend only on the actions of the players in the other partition. More specifically, for player $i \in \mathcal{G}_j$, $j \in \{0, 1\}$, and for $p \in \Delta$,

$$\bar{u}_i(p) = g_i(p^{-j}),$$

for some function $g_i : \left(\prod_{k \in \mathcal{G}_{(j+1) \bmod 2}} S_k \right) \rightarrow \mathbb{R}^{J_i}$, where p^{-j} is the mixed strategy profile of all players that are not in partition \mathcal{G}_j . Therefore, for $x, y \in S$,

$$\sigma_{iy_i}(\bar{u}_i(x)) = \sigma_{iy_i}(g_i(x^{-j})).$$

Therefore,

$$p_{xy} = \prod_{i \in \mathcal{G}_0} \sigma_{iy_i}(g_i(x^1)) \times \prod_{i \in \mathcal{G}_1} \sigma_{iy_i}(g_i(x^0)),$$

where x^0 and x^1 are the mixed strategy profiles of all the players in partition \mathcal{G}_0 and \mathcal{G}_1 , respectively.

Lemma 1. \mathcal{P} maps product measures to product measures.

Proof. Let $\mathcal{S}_0 = \prod_{k \in \mathcal{G}_0} S_k$ and $\mathcal{S}_1 = \prod_{k \in \mathcal{G}_1} S_k$ denote the sets of strategy profiles restricted to \mathcal{G}_0 and \mathcal{G}_1 . Suppose $\mu \in \Delta$ can be decomposed as $\mu_0 \times \mu_1$, where μ_0 is a measure on \mathcal{S}_0 and μ_1 is a measure on \mathcal{S}_1 . Then

$$\begin{aligned} \mu \mathcal{P} &= \left(\sum_{x \in S} \mu_1(x^1) \prod_{i \in \mathcal{G}_0} \sigma_{iy_i}(g_i(x^1)) \times \mu_0(x^0) \prod_{i \in \mathcal{G}_1} \sigma_{iy_i}(g_i(x^0)) \right)_{y \in S} \\ &= \left(\left[\sum_{x^1 \in \mathcal{S}_1} \mu_1(x^1) \prod_{i \in \mathcal{G}_0} \sigma_{iy_i}(g_i(x^1)) \right] \times \left[\sum_{x^0 \in \mathcal{S}_0} \mu_0(x^0) \prod_{i \in \mathcal{G}_1} \sigma_{iy_i}(g_i(x^0)) \right] \right)_{y \in S} \\ &= \nu_0 \times \nu_1, \end{aligned}$$

where

$$\begin{aligned}\nu_0(y^0) &= \sum_{x^1 \in \mathcal{S}_1} \mu_1(x^1) \prod_{i \in \mathcal{G}_0} \sigma_{iy_i^0}(g_i(x^1)), \\ \nu_1(y^1) &= \sum_{x^0 \in \mathcal{S}_0} \mu_0(x^0) \prod_{i \in \mathcal{G}_1} \sigma_{iy_i^1}(g_i(x^0))\end{aligned}$$

as desired. \square

Lemma 2. *The set of product probability measures is sequentially closed.*

Proof. Let $X = X_0 \times X_1$, and Π denote the set of corresponding product probability measures on X , and fix a sequence $(\mu_n)_n = (\mu_{0_n})_n \times (\mu_{1_n})_n \in \Pi$ with limit μ . For $A \subset X_0, B \subset X_1$, $\mu(A \times B) = \lim_n \mu_n(A \times B) = \lim_n (\mu_{0_n}(A) \mu_{1_n}(B)) = \lim_n (\mu_n(A \times X_1) \mu_n(X_0 \times B)) = \lim_n (\mu_n(A \times X_1)) \cdot \lim_n (\mu_n(X_0 \times B)) = \lim_n (\mu_{0_n}(A)) \cdot \lim_n (\mu_{1_n}(B)) \in \Pi$. \square

Since \mathcal{P} is the transition matrix of an ergodic Markov chain, there is a unique equilibrium distribution π , such that for all initial distributions λ , $\pi = \lim_n \lambda \mathcal{P}^n$. Let λ_0 denote a measure in Δ that can be decomposed as a product of measures on \mathcal{S}_0 and \mathcal{S}_1 (such as any pure strategy profile). By the above lemmas, $\pi = \lim_n \lambda_0 \mathcal{P}^n$ is a product measure.

References

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